

Message Queue Formats

How does the message queue work?

Sep 14, 1988

This note describes the format of the shared memory Message Queue used for communication with another processor on the VMEbus. It is implemented as a simple one-way queue. Messages are placed in the queue by the VME System computer. The other processor removes messages from the queue and interprets the command accordingly.

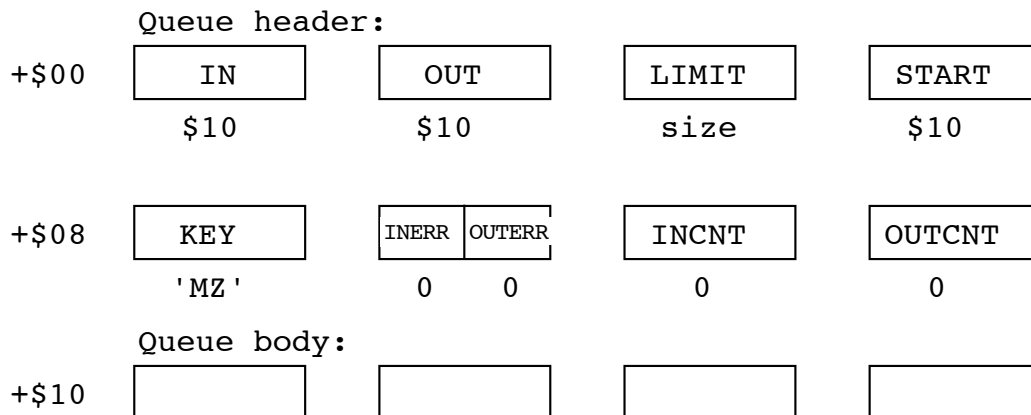
Initialization

The VME System computer initializes co-processor message queues at reset time. One of the standard system tables—table #15—contains pointers to the message queue for each co-processor. This queue pointer table, indexed by co-processor number, has 8-byte entries of the following format:



The message queue pointer is followed by the total queue size.

Each co-processor queue has the following format:



The values beneath the words in the queue header are the initialized values. The other processor, when it recognizes the presence of the queue, examines the KEY field. If it has the value 'MZ' it changes it to 'MQ' to signal to the VME system cpu that it has “seen” the message queue. (Until this happens, the VME system cpu will not place messages into the queue.)

Queue header

The IN pointer (offset from the start of the queue header) points to the next available space in the queue for a message. It is altered by the VME system cpu as the last act upon placing the new message into the queue, after first checking that there is room available to hold the message.

The OUT pointer points to the next message to be removed from the queue by the co-processor. It is altered after the co-processor has removed the message from the queue. When the two pointer IN and OUT are equal, the queue is considered empty. When they are unequal, there is at least one message in the queue.

The **LIMIT** word is the total size of the queue (in bytes). It is determined by the contents of the VME system table directory.

The **START** word is the offset to the start of the queue body. When new entries have reached **LIMIT**, the **IN** pointer circles back to **START**.

The **KEY** word has the value 'MZ' when the queue is initialized, and it is changed by the co-processor to 'MQ' to signal that the queue has been recognized.

The **INERR** byte counts times when the VME system cpu tried to place an entry into the queue, but found the queue full.

The **OUTERR** byte is incremented by the co-processor cpu when it encounters an error in processing the messages it removes from the queue.

The **INCNT** word is incremented for each message successfully placed into the queue.

The **OUTCNT** word is incremented by the co-processor when it successfully removes a message from the queue.

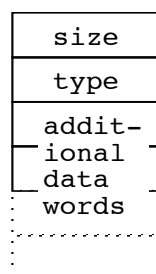
Protocol

When the VME system cpu has a message to place into the queue, it checks to see that the **KEY** word has the value 'MQ'. It then checks to see if the message can fit. If $IN \geq OUT$, it checks for space between **IN** and **LIMIT**; if there is not enough space there, a zero is placed at the word pointed to by **IN**, and **IN** is reset to **START**. If either $IN < OUT$ or **IN** had to be reset to **START**, it checks for space between **IN** and **OUT**. If there is room, the message is copied into the space, and the **IN** pointer is advanced by the message size.

The co-processor examines the queue at its convenience. (Note that the co-processor had to have *a priori* knowledge of the location of the queue.) If the queue is not empty ($IN \neq OUT$), a message is removed from the queue. The word pointed to by **OUT** is examined. If it is zero, **OUT** is reset to **START**, and if $IN \neq OUT$, the word at **START** is examined. When the co-processor has removed the message from the queue, it advances the **OUT** pointer by the message size. It may then check to see if another message is present. It is assumed that the co-processor will poll the message queue often enough that the queue will not fill up. If it does, one will find the **INERR** count nonzero.

Message format

Messages placed in the queue for a co-processor conform to a simple structure:



The first word is the message size, including the `size` word. The second word is the message `type`. Additional message contents may follow the second word. So the minimum message size is 4 bytes, in the case that no additional data is required for a given type. The `type` word and any additional data have meaning only to the co-processor, not the VME system.

Generic message setting

To send a general message to a co-processor in a VME system, one has only to send the appropriate setting, specifying listype #40. The ident used with this listype supplies the co-processor number (0,1,2...). The number of data bytes specified in the setting—incremented by 2—becomes the `size` word of the message placed into the selected queue. The first word of the setting data, then, is the `type` word of the message. Additional data words follow the `type` word. Note that the VME system does not care about the `type` word value.

Analog control

There is a format in use for analog control, in which an analog channel may be set which results in a message sent to a co-processor queue.

size	
\$00	type
index	
data	

In this case, the `type` value is given by a byte from the analog control field of the analog descriptor for that channel. The `index` value is given by a word from the analog control field. The data word is the word of setting data in the analog control setting. The `size` value is therefore 8 bytes.